Optimal Partition for Parallel Rendering

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MAR.06.2008
Optimal Partition for Parallel Rendering

• Load balanced partition for rendering
  – Adhoc and Scan-like Tries
  – No Optimal Approaches for reducing Parallel rendering Hazards

• Partition algorithm for Parallel Rendering
  – Load Balance with Minimization Parallel Rendering Hazards
  – Realtime Partitioning Performance
# Parallel Graphics Hazards

## Processing Hazard
- Unbalanced Graphics Processing
  - vertices and pixel processing
  - overlap primitives between different partitions

## Management Hazard
- Unbalanced Management Processing
  - frame coherency between each processor’s partitions of animate frame sequence
  - out of core rendering resources management
Optimal Partition Generation for Parallel Graphics

Typical Load balancing for Parallel processing
- Preprocessed, Static Partition
- Load Balance with Minimization Cutting Edge between Partitions (for communication minimization)

New Load balancing for Parallel Realtime Graphics
- Runtime, Dynamic Partition with computing time constraints
- Load Balance with minimizing Parallel Graphics Hazards
Optimal Partition Generation for Parallel Graphics

- **Visibility Culling**
  - View frustum Culling with Hierarchical Structure
  - Problem set per each image frame

- **Graph Modeling**
  - Adapting graph size for realtime partition algorithm
  - 3D or 2D Graph for problem set
  - Graphics Hazards Modeling via Weighted Laplacian

- **Graph Partitioning**
  - Hazards Minimization within load balance condition
  - Fast, Efficient and Scalable Graph Partition Algorithm with realtime constraints
  - Optimal graph partition for Assignments to graphic processors
Graph Partition Algorithm - Overview of Spectral Bisection

- A Workload Graph expresses as a Laplacian Matrix.
- Rayleigh quotient of Assignment Vector and laplacian matrix represents sum of cutting edge cost between two partitions.

<table>
<thead>
<tr>
<th>Laplacian Matrix</th>
<th>$L_{ij} = \begin{cases} -1 &amp; \text{if } (v_i, v_j) \in E \ d_i &amp; \text{if } i = j \ 0 &amp; \text{else} \end{cases}$</th>
</tr>
</thead>
<tbody>
<tr>
<td>Discrete Optimization</td>
<td>$x^t Lx = 4 \times {# \text{ edges connecting nodes in N - to nodes in N +}}$</td>
</tr>
<tr>
<td></td>
<td>$x_i = \pm 1$</td>
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<tr>
<td>Continues Relaxation</td>
<td>$F(x) = \frac{x^t Lx}{x \cdot x} = \sum_{(i,j) \in E} (x_i - x_j)^2$</td>
</tr>
</tbody>
</table>
Graph Partition Algorithm - Overview of Spectral Bisection

- Spectral analysis of laplacian matrix; Second smallest eigenvector minimizes Rayleigh Quotient.
- Optimal assignment vector is a good approximation to load balanced graph partition problem with minimal communication edges between partitions.

### Spectral Analysis

\[
L = \sum_{i=1}^{n} \lambda_i z_i z_i^t \quad \lambda_1 = 0, z_1 = e \text{ and } \lambda_1 < \lambda_{i+1}
\]

### Optimal Assignment Vector

\[
\text{minimize } \frac{x^t L x}{x \cdot x} \text{ over } x \Rightarrow x = z_2, \quad \frac{x^t L x}{x \cdot x} = \lambda_2
\]
Graph Partition Algorithm – Algorithm Selection

Avoidance of computing full spectral analysis; just need is second smallest Eigen vector.

Modification of conventional optimization method;
Conjugate Gradient Method

- Shows High convergence speed to second smallest Eigen vector of Laplacian matrix
- Meets Load balance constraint
- Requires Low memory
Conjugate gradient method for the spectral partitioning of graphs

\[ F(x) = \frac{x^t L x}{x \cdot x}, \quad \nabla F(x) = \frac{2}{x \cdot x} (L \cdot x - F(x)x) \]

Step 1: Initialize \( x_0 \) and gradient vector \( g_0 = -\nabla F(x_0), h_0 = g_0 \)

Step 2: Fletch-Reeves Conjugate gradient method

\[ \alpha_k = \min_{\alpha} F(x_{k-1} + \alpha h_{k-1}) \]
\[ x_k = x_{k-1} + \alpha_k h_{k-1} \]
\[ g_k = -\nabla F(x_k) \]
\[ h_k = g_k + \frac{g_k \cdot g_k}{g_{k-1} \cdot g_{k-1}} h_{k-1} \]

Step 3: Check for convergence.
If convergence has not been reached go to Step 2.

Step 4: Accept optimal vector \( x = x_k \)
Experiment for feasibility study

Visibility Culling
- View frustum Culling with Hierarchical Structure
- Visible Octree cell set per each image frame

Graph Modeling
- 3D 6 Grid Vertex Graph for problem set
- Rendering Primitive Cost Modeling via Weighted Laplacian

Graph Partitioning
- User Specific Hazard and Load balance thresholds
- Fast, Efficient and Scalable Graph Partition Algorithm with realtime constraints
- Optimal graph partition for Assignments to graphic processors
Load Balanced Partition Generation

Urban Scene (Digital Media City)
323863 triangles
[262 visible octree cells]

Positive Set
164524
[49]

Negative Set
159339
[213]

Imbalance
5185 / 323863 = 1.60%
Optimal Bisection for a image frame

182 iteration for partitioning weighted graph of 262 octree cells
Optimal Bisection for animation frames

Camera Closing to urban ground and Partitions
Dynamic Workload by camera movement

Graph Size and Imbalance

With Objective function threshold 0.0002, Imbalance threshold 2%
Graph Size and Iteration Count

Graph Size and Execution Time

VertexCount  IterCount

Intel Pentium 2.4Ghz
3Giga Memory
Conclusion

• Optimal partition could be solved in real-time for parallel graphics and simulations.

• Future works
  – Graphics Hazards Modeling
    • for reducing parallel graphics hazards
  – Valley Constraint Optimization
    • for partition’s frame coherency
  – Recursive partitioning
    • for K-way processors
Thanks a lot.